

Optimizing Hierarchy Joins

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Real Life Scenario

- Brewery Inc.
 - Daily reports for production and sales
 - Need to organize data in data warehouse
- Options:
 - Single table per production/sales
 - Table per production/sales per day



Options Discussion

- Single table per production/sales
 - + Simple queries for yearly/quarter reports
 - Performance issues with daily reports
- Table per production/sales per day
 - + Good performance for daily reports
 - Complicated queries for yearly/quarter reports
- Want: benefits from both



PostgreSQL Support

- Inheritance
 - Way to organize tables
 - Retain the abstraction of a single table
- Example

production (id, date, quantity, ...)
production_jan_1 (id, date, quantity, ...)
production_jan_2 (id, date, quantity, ...)

- Query over "production" => combines all data
- Query over "production_jan_1" => only Jan 1st



Specifics

• Parent table

create table production (int "id", timestamp "date", ...);

Child tables

create table production_jan_1 (
 check ('2009-01-01' <= "date" and "date" < '2009-01-02')
) inherits (production);</pre>



Sample Reports

- Beer quantities produced on Jan 1st select id, sum(quantity) from production where '2009-01-01' <= "date" and "date" < '2009-01-02' group by id;
- Execution plan scans ONLY "production_jan_1"
- How?
 - Filters tables based on check constraints



Sample reports cont.

- Sales per beer kind per day select p.id, p.date, sum(s.price) from production p, sales s where p.date = s.date and p.id = s.id group by p.id, p.date;
- Execution plan treats each hierarchy as single table
 - next slide



Execution Plan

Group agrregate

- -> merge join
 - -> sort
 - -> Append
 - -> Seq Scan on production
 - -> Seq Scan on production_jan_1

-> . . .

- -> sort
 - -> Append
 - -> Seq Scan on sales
 - -> Seq Scan on sales_jan_1

-> . . .



Execution Plan - Discussion

- Extra work
 - matching beer produced on Jan 1 with sales from Jan 2
- Can we avoid it?
 - Utilize check constraints
 - Join child tables directly



New Execution Plan

Hash aggregate

-> Append

-> . . .

- -> Hash Join
 - -> Seq Scan on production_jan_1
 - -> Seq Scan on sales_jan_1
- -> Hash Join
 - -> Seq Scan on production_jan_2
 - -> Seq Scan on sales_jan_2

ASTER

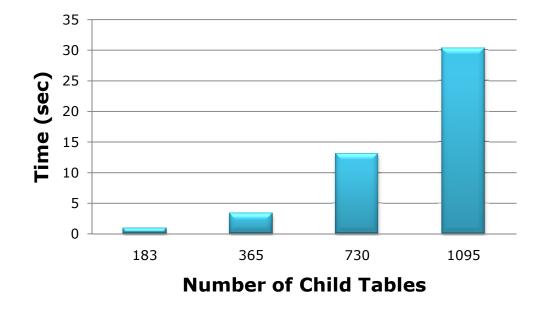
Algorithm Overview

- 1. Check if we can join child tables directly
- 2. Get child table constraints
- 3. Find possible child joins
- 4. Generate plans for each child join
- 5. Combine results from child joins



Find possible child joins

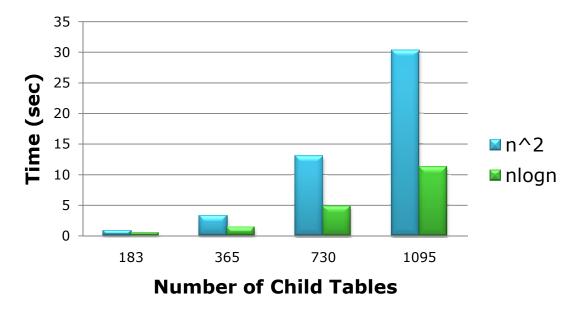
- First implementation: naïve n^2
 - Examine constraints from each possible pair of child tables
 - Expensive





Find possible child joins cont.

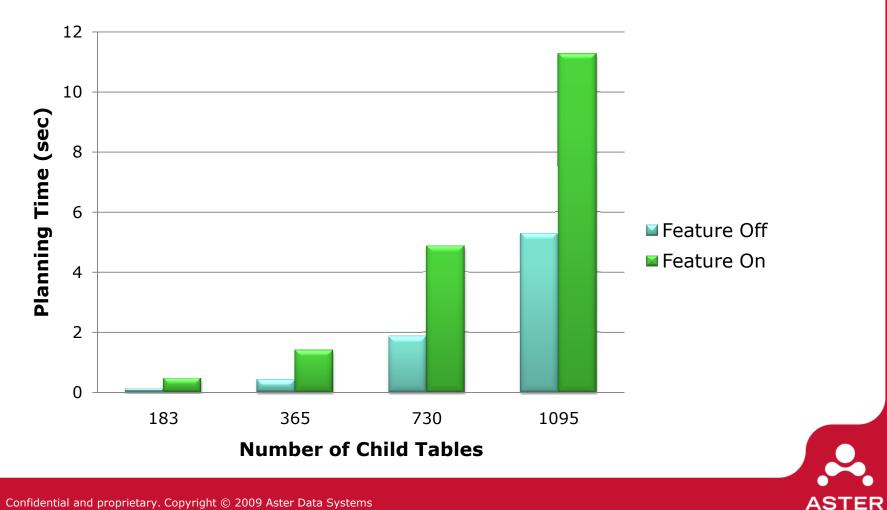
- Another approach:
 - Treat constraints as intervals
 - Use an interval tree for the matching of child tables
 - Complexity: n*logn



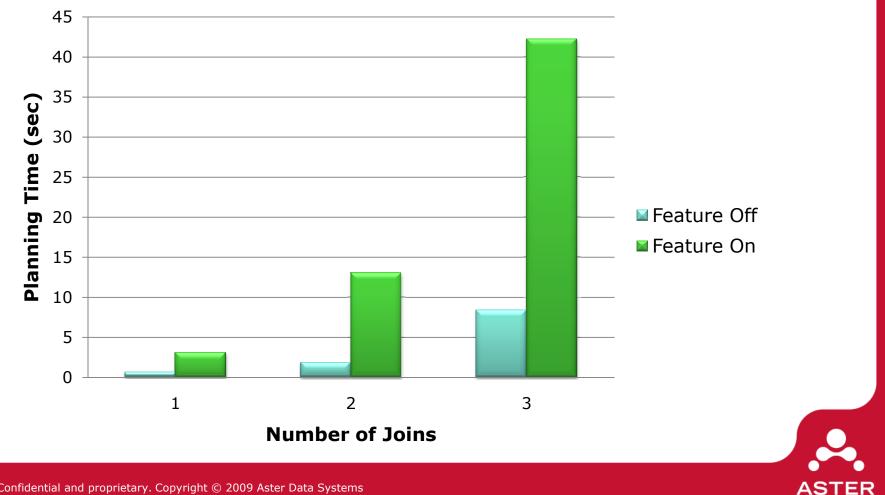
- Goals:
 - Examine overhead on planning time
 - Examine effects on execution time
- Factors
 - Number of child tables
 - Number of records in hierarchies
 - Number of joins in query



• Number of joins = 2



• Number of child tables = 730



- Number of child tables = 365
 - 30 (sec) 25 **Total Execution Time** 20 15 Feature Off Feature On 10 5 0 0.1 0.2 0.3 0.4 0.5 Number of Tuples Per Hierarchy (million)

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• Number of joins = 2